# LEMP: Fast Retrieval of Large Entries in a Matrix Product 

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## Recommender Systems

- Problem
- Set of users
- Set of items (movies, books, jokes, products, stories, ...)
- Feedback (ratings, purchase, click-through, tags, ...)
- Predict the preference of each user for each item



## Latent Factor Models



## Recommender Systems

- Given
- A query matrix $\mathbf{Q}$



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- A probe matrix $\mathbf{P}$

Adam
Bob
Charlie
Dennis $\left(\begin{array}{cc}3.2 & -0.4 \\ 3.1 & -0.2 \\ 0 & 1.8 \\ -0.4 & 1.9 \\ \mathbf{Q}^{T}\end{array}\right)$


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- Given
- A query matrix $\mathbf{Q}$
- A probe matrix $\mathbf{P}$
- A threshold $\theta>0$
- Find good recommendations

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- Find good recommendations
- All entries in $\mathbf{Q}^{T} \mathbf{P}$ that are $\geq \theta$
- Each entry is an inner product $\mathbf{q}^{T} \mathbf{p}=\sum_{i=1}^{r} q_{i} p_{i}$

|  |  |  |  | $00^{\text {a }}$ |  | 人i |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $\left(\begin{array}{l}1.6 \\ 0.6\end{array}\right.$ | 1.3 0.8 | 0.7 2.7 |  | 0.4 2.2 |
| Adam | $\left(\begin{array}{ll}3.2 & -0.4 \\ 3.4\end{array}\right.$ | 4.9 | 3.8 | 1.2 | 2.1 | 0.4 |
| Bob | $\left(\begin{array}{ll}3.1 & -0.2\end{array}\right.$ | 4.8 | 3.9 | 1.6 | 2.5 | 0.8 |
| Charlie | $\left(\begin{array}{ll}0 & 1.8\end{array}\right.$ | 1 | 1.4 | 4.9 | 5.0 | 4.0 |
| Dennis | $\left(\begin{array}{cc}-0.4 & 1.9\end{array}\right)$ | 0.5 | 1 | 4.9 | 4.9 | 4.0 |
|  | $\mathbf{Q}^{\text {T }}$ |  |  | $\mathbf{Q}^{T} \mathbf{P}$ |  |  |

## Problem Statement

Maximum Inner Product Search
Find pairs of vectors with large inner products
Given

- a query vector $\mathbf{q}^{T}$ (from matrix $\mathbf{Q}^{T}$ )
- a set $\mathbf{P}$ of probe vectors (the matrix $\mathbf{P}$ )
- a threshold $\theta>0$

Find

- all vectors $\mathbf{p}$ such that $\mathbf{q}^{\top} \mathbf{p} \geq \theta$


## Naive Solution

- Compute full matrix product $\mathbf{Q}^{T} \mathbf{P}$
- Determine which entries are $\geq \theta$
- Complexity $O$ (mnr)
- Usually
- $\mathrm{m}, \mathrm{n}$ : order of millions
- $10<r<500$
- Example
- $\mathrm{m}=10$ millions
- $\mathrm{n}=1$ million
- \#entries = 10 trillion
- Avg. Inner Product Time $=100$ nsec

- Runtime > 11 days
- Can we do better than that?

When is an inner product large?

$$
\mathbf{q}^{\top} \mathbf{p}=\|\mathbf{q}\|\|\mathbf{p}\| \cos \angle(\mathbf{q}, \mathbf{p}) \geq \theta,-1 \leq \cos \angle(\mathbf{q}, \mathbf{p}) \leq 1
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| Length <br> $(\\|\mathbf{q}\\|\\|\mathbf{p}\\|)$ | short <br> $(<\theta)$ | medium <br> $(\approx \theta)$ |
| :--- | :--- | :--- |
| Angle <br> ( $\cos \angle(\mathbf{q}, \mathbf{p}))$ | small positive <br> large positive |  |
| Suitable <br> method |  |  |

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| :--- | :--- | :--- | :--- |
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| Suitable pall positive <br> large positive <br> method | Prune |  |  |

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Cosine Similarity
Search $\theta^{\prime}=\frac{\theta}{\|\mathbf{q}\|\|\mathbf{p}\|}$
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$\xrightarrow[(\|\mathbf{q}\|\|\mathbf{p}\|)]{\text { Length }}$

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Naive-like retrieval

## Main idea: bucketize by length

- Partition $\mathbf{P}$ in buckets with vectors of similar length
- Index buckets suitably



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- Partition $\mathbf{P}$ in buckets with vectors of similar length
- Index buckets suitably
- For each query vector and bucket
- Determine local threshold
- Prune bucket if possible


Local threshold on cosine similarity:

$$
\theta_{b}(\mathbf{q})=\frac{\theta}{\|\mathbf{q}\| l_{b}}
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## Main idea: bucketize by length

- Partition $\mathbf{P}$ in buckets with vectors of similar length
- Index buckets suitably
- For each query vector and bucket
- Determine local threshold
- Prune bucket if possible
- Otherwise, select best retrieval method


Local threshold on cosine similarity:

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\theta_{b}(\mathbf{q})=\frac{\theta}{\|\mathbf{q}\|_{b}}
$$

## Discussion

LEMP

- Prunes whole buckets
- Selects a suitable retrieval algorithm per query and bucket
- Can leverage existing methods
- Cache-friendly


## Bucket-level retrieval

Choose among a variety of algorithms

- Fagin's Threshold Algorithm ( $r<10$ )
- All pairs similarity search family $(1000<r)$
- Space-partitioning trees $(r<10)$

Our vectors are

- Not necessarily sparse
- Real values
- Medium dimensionality ( $10<r<500$ )

Two new algorithms

- COORD
- INCR


## INCR: Main Idea

- $\mathbf{q}, \mathbf{p}$ qualify if $\|\mathbf{q}\|\|\mathbf{p}\|\left(\sum_{i=1}^{r} \bar{q}_{i} \bar{p}_{i}\right) \geq \theta$

| $\\|\mathbf{q}\\|$ | $\overline{\mathbf{q}}_{1}$ | $\overline{\mathbf{q}}_{2}$ | $\overline{\mathbf{q}}_{3}$ | $\overline{\mathbf{q}}_{4}$ | $\ldots$ | $\overline{\mathbf{q}}_{r}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
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- In practice
- $\phi$ is automatically tuned
- the $\phi$ coordinates do not have to be consecutive

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INCR
$\square$

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$\square$

id $\sum$ UB


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## INCR

- Pick coordinates (largest $\left|\bar{q}_{i}\right|$ first)
- Scan indexes and update " $\sum$ " and "UB" quantities
- Prune vectors for which $\|\mathbf{q}\|\|\mathbf{p}\|\left(\sum+U B\right)<\theta$
- No need to scan the whole lists! Bounds exist



## Discussion

INCR

- Computes partial inner products
- Uses simple index - cheap to construct
- Scans part of the index - does not necessarily scan from top
- Sequential memory access pattern, fast


## How fast is it? (Top-1 movie per user)



## Summary

## Large inner-product search

- Matrix factorization common technique in data mining
- Large entries in matrix products are usually of particular interest

The LEMP algorithm

- Bucketizes vectors by length
- Prunes buckets whenever possible
- For the remaining buckets: selects efficient retrieval algorithms
- Consistently fastest bucket method: INCR


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## Thank you! <br> Questions?

## Performance of bucket algorithms



IE-SVD ${ }^{T}$


## Do we need to scan the whole lists?



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## COORD: bounds

- Find for each coordinate $\bar{p}_{f}$ of $\overline{\mathbf{p}}$ an upper and lower bound $\left[L_{f}, U_{f}\right]$ such that $\bar{p}_{f} \notin\left[L_{f}, U_{f}\right] \Rightarrow \overline{\mathbf{q}}^{\top} \overline{\mathbf{p}}<\theta_{b}(\mathbf{q})$



## Algorithm selection



