Database Technology
Entity Relationship Models
Previously on Database Technology

• Introduction to Relational Databases
  – A standard model for storing data
  – Using relations/tables

• Introduction to SQL
  – Creating and changing tables
  – Reading and writing data into tables
Today

- Designing databases
  - i.e., how to get from your customer’s requirements...
  - ...to a set of tables and attributes
Outline

• Design Process
• Modeling
• Constraints
• E-R Diagrams
• Design Issues
• Weak Entity Sets
• Extended E-R Features
• Reduction to Relation Schemas
• Comparison UML
Database Design

• Initial phase: requirements engineering
  – characterize fully the data needs of the prospective database users
  – which data needs to be stored?
    • ...and in which volumes?
  – which queries should be answered?

• Conceptual schema
  – which types of entities and relations exist?
  – what attributes do they have?
Database Design

• Final phase: from a conceptual to physical data model
  – Logical Design: find a “good” collection of relation schemas
    • Business decision – What attributes should we record in the database?
    • Computer Science decision – What relation schemas should we have and how should the attributes be distributed among the various relation schemas?
  – Physical Design – Deciding on the physical layout of the database
Database Design Approaches

• Entity Relationship Model (today)
  – Models an enterprise as a collection of entities and relationships
  – Entity: a “thing” or “object” in the enterprise that is distinguishable from other objects
    • Described by a set of attributes
  – Relationship: an association among several entities
  – Represented diagrammatically by an entity-relationship diagram

• Normalization Theory (next lecture)
  – Formalize what designs are bad, and test for them
Entity Relationship Model

• Dates back to the 1970s
  – developed to facilitate database design by allowing the specification of an enterprise schema that represents the overall logical structure of a database

• Toolkit for mapping the meanings and interactions of real-world enterprises onto a conceptual schema

• The ER data model employs three basic concepts:
  – entity sets,
  – relationship sets,
  – attributes

• Associated diagrammatic representation (ER diagram)
  – graphic expression of the overall logical structure of a database
Entity Sets

• An entity is an object that exists and is distinguishable from other objects
  – Example: Peter Chen, Mannheim, Star Wars

• An entity set is a set of entities of the same type that share the same properties
  – Example: set of all persons, cities, movies

• Each entity is represented by a set of attributes
  – Example:
    
    instructor = (ID, name, street, city, salary )
    
    course= (course_id, title, credits)

• A subset of the attributes form a primary key of the entity set
  i.e., uniquely identifying each member of the set
Entity Sets – Example

- instructor (instructor_id, instructor_name)
- student (student_id, student_name)

<table>
<thead>
<tr>
<th>instructor_id</th>
<th>Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>76766</td>
<td>Crick</td>
</tr>
<tr>
<td>45565</td>
<td>Katz</td>
</tr>
<tr>
<td>10101</td>
<td>Srinivasan</td>
</tr>
<tr>
<td>98345</td>
<td>Kim</td>
</tr>
<tr>
<td>76543</td>
<td>Singh</td>
</tr>
<tr>
<td>22222</td>
<td>Einstein</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>student_id</th>
<th>Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>98988</td>
<td>Tanaka</td>
</tr>
<tr>
<td>12345</td>
<td>Shankar</td>
</tr>
<tr>
<td>00128</td>
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</tr>
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<td>76543</td>
<td>Brown</td>
</tr>
<tr>
<td>76653</td>
<td>Aoi</td>
</tr>
<tr>
<td>23121</td>
<td>Chavez</td>
</tr>
<tr>
<td>44553</td>
<td>Peltier</td>
</tr>
</tbody>
</table>
Relationship Sets

- A **relationship** is an association among several entities
  
  Example:
  
  44553 (Peltier) **advisor** 22222 (Einstein)
  
  *student* entity  **relationship set**  *instructor* entity

- A **relationship set** is a mathematical relation among \( n \geq 2 \) entities, each taken from entity sets

  \[
  \{(e_1, e_2, \ldots, e_n) \mid e_1 \in E_1, e_2 \in E_2, \ldots, e_n \in E_n\}
  \]
  where \((e_1, e_2, \ldots, e_n)\) is a relationship

- Example:

  \((44553, 22222) \in advisor\)
Relationship Sets

```plaintext
<table>
<thead>
<tr>
<th>instructor</th>
<th>student</th>
</tr>
</thead>
<tbody>
<tr>
<td>76766</td>
<td>98988</td>
</tr>
<tr>
<td>45565</td>
<td>12345</td>
</tr>
<tr>
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<td>98345</td>
<td>76543</td>
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<tr>
<td>76543</td>
<td>76653</td>
</tr>
<tr>
<td>22222</td>
<td>23121</td>
</tr>
<tr>
<td></td>
<td>44553</td>
</tr>
<tr>
<td>Crick</td>
<td>Tanaka</td>
</tr>
<tr>
<td>Katz</td>
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<td>Singh</td>
<td>Aoi</td>
</tr>
<tr>
<td>Einstein</td>
<td>Chavez</td>
</tr>
<tr>
<td></td>
<td>Peltier</td>
</tr>
</tbody>
</table>
```
Relationship Sets

- An attribute can also be associated with a relationship set
- E.g., *advisor* relationship set:
  - date which captures the start of the supervision
Degree of a Relationship

- Definition: degree of a relationship
  - i.e., number of entity sets that are involved in relation set
- Binary relationship (degree two)
  - involve two entity sets
  - the by far most frequent case
- Relationships between more than two entity sets (degree >2)
  - e.g.: students work on projects under the guidance of an instructor
  - relationship proj_guide is a ternary relationship
    between instructor, student, and project
  - those are rather rare
Cardinality Constraints

• Express the number of entities to which another entity can be associated via a relationship set
  – Most useful in describing binary relationship sets

• For a binary relationship set, the mapping cardinality must be one of the following types:
  – 1:1 (one to one)
  – 1:n (one to many)
  – n:1 (many to one)
  – n:m (many to many)
Mapping Cardinalities – One to One

• One to one (1:1)
  – Note: Some elements in A and B may not be mapped to any elements in the other set

• Examples
  – student_works_on_thesis
  – department_has_dean
Mapping Cardinalities – One to Many

• One to many (1:n)
  – Note: Some elements in A and B may not be mapped to any elements in the other set

• Examples
  – building_has_room
  – course_taught_by_lecturer
Mapping Cardinalities – Many to One

• Many to one (n:1)
  – Note: Some elements in A and B may not be mapped to any elements in the other set

• Examples
  – room_located_in_building
  – lecturer_teaches_course

• Note:
  – the inverse of a 1:n relation is a n:1 relation
  – and vice versa
Mapping Cardinalities – Many to Many

- Many to many (n:m)
  - Note: Some elements in A and B may not be mapped to any elements in the other set

- Examples
  - student_takes_course
  - student_has_advisor
Distinguishing 1:n/n:1 and n:m Cardinalities

- **Rule of thumb**
  - Always ask for the cardinality the other way around

- “A building may have multiple rooms...”
  - “...but can a room be in multiple buildings?”
  - No → building_has_room is 1:n

- “A department can be located in multiple buildings...”
  - “...but can a building host multiple departments?”
  - Yes → department_located_in_building is n:m
Relation Sets from the Same Entity Set

• The two entity sets in a relation set may be the same
• This holds independently from the cardinality!

- person_married_to_person
  - 1:1
- person_is_father_of_person
  - 1:n
- person_has_father
  - n:1
- person_is_parent_of_person
  - n:m
Attribute Types & Domains

• Attribute types:
  – Simple and composite attributes
  – Single-valued and multi-valued attributes
    • Example: multi-valued attribute: phone_numbers
• Derived attributes
  – Can be computed from other attributes
    – Example: age (given date_of_birth)
• Domain – the set of permitted values for each attribute
Composite Attributes

Composite attributes

name

first_name middle_initial last_name

address

street city state postal_code

street_number street_name apartment_number

component attributes
Redundant Attributes

• Suppose we have entity sets:
  – instructor, with attributes: ID, name, dept_name, salary
  – department, with attributes: dept_name, building, budget

• In ERM, instructors and departments are connected by a relation set
  – e.g., instructor_belong_to_department (ID,dept_name)

• Now, dept_name is no longer needed in the instructor entity set
  – It is redundant there
  – Hence, we will remove it

• *Note*: sometimes, removed redundant attributes are reintroduced when converting the conceptual model into a logical model
Weak Entity Sets

• Consider the set of buildings and rooms
  – Entity set building(building_name,address)
  – Entity set room(number,capacity)
  – Relation set room_in_building (number,building_name)

• Note:
  – As in the previous example, we have removed the redundant attribute
    building_name from the entity set room

• Question:
  – What is the primary key of the the entity set room?
Weak Entity Sets

• Weak entity sets are entity sets that
  – do no not have a set of attributes sufficient to identify each entity uniquely
  – require an additional relation set to identify each entity uniquely
• Those relation sets are called *identifying relation set*

• Weak entities do not have primary keys
  – A weak entity set has an *identifying entity* and a *discriminator*
  – Example:
    • *building* is the identifying entity
    • *number* is the discriminator

• A weak entity cannot exist without the identifying entity
  – e.g., a room cannot exist without the building
**ER Diagrams**

- **Entity Relationship Diagrams (ER diagrams)**
  - are the graphical notation of entity relationship models

- **Notation of entity sets**:
  - Rectangles represent entity sets
  - Attributes listed inside entity rectangle
  - Underlining indicates primary key attributes

![ER Diagram Examples](image-url)
ER Diagrams

- Diamonds represent relationship sets
ER Diagrams

- Diamonds represent relationship sets
  - Attributes can be attached to relationship sets

```
+----------------+       +----------------+
| instructor     |       | student         |
| ID             |       | ID              |
| name           |       | name            |
| salary         |       | tot_cred        |
+----------------+       +----------------+
                  \   |                |
                  date |
+----------------+       +----------------+
```


Roles

- Entity sets of a relationship need not be distinct
  - i.e., there may be a relationship set involving the same entity set twice
- Each occurrence of an entity set plays a “role” in the relationship
  - The labels “course_id” and “prereq_id” are called roles

```
course
<table>
<thead>
<tr>
<th>course_id</th>
<th>title</th>
<th>credits</th>
</tr>
</thead>
</table>

prereq
| course_id | prereq_id |
```
Cardinalities in ER Diagrams

- We express cardinality constraints by drawing either a directed line (→), signifying “one”, or an undirected line (—), signifying “many”, between the relationship set and the entity set.

- One-to-one relationship between an instructor and a student:
  - A student is associated with at most one instructor via the relationship advisor
  - An instructor is associated with at most one student via the relationship advisor
Cardinalities in ER Diagrams

- one-to-many relationship between an instructor and a student
  - an instructor is associated with several (including 0) students via advisor
  - a student is associated with at most one instructor via advisor
Cardinalities in ER Diagrams

- Many to many relationships
  - An instructor is associated with several (possibly 0) students via advisor
  - A student is associated with several (possibly 0) instructors via advisor
Total and Partial Participation

- Total participation (double line)
  - every entity in the entity set participates in at least one relationship in the relationship set
  - i.e., every student *must* have an advisor
    - recap: think of *not null* constraints

- Partial participation (single line)
  - some entities may not participate in the relationship
  - e.g., not every instructor has to supervise a student
Complex Cardinality Constraints

- Notation for minimum/maximum cardinality of a relation
  - Each student has *exactly one* advisor (i.e., \( \text{min} = \text{max} = 1 \))
  - Each instructor can be the advisor of multiple students, but needs not be (i.e., \( \text{min} = 0, \text{max} = \infty \))

- Notation:
  - \( \text{min..max} \)
  - * indicates no limit
Notation of Attribute Types

instructor

ID
name
  first_name
  middle_initial
  last_name
address
  street
    street_number
    street_name
    apt_number
city
state
zip
  \{ phone_number \}
date_of_birth
age ()

complex attribute

multivalued attribute

derived attribute
Expressing Weak Entity Sets

- A weak entity set is depicted via a double rectangle
  - The identifying relationship set is depicted by a double diamond
- The *discriminator* is underlined with a dashed line
  - Primary key for section – (course_id, sec_id, semester, year)
Higher Arity Relationship Sets

- Most relationship sets are binary
- Sometimes, ternary (or higher arity) relations occur
  - ER models support that
- Example:
  - Students work on projects under supervision of an instructor

```
project

instructor
ID
name
salary

student
ID
name
tot_cred
```

proj_guide

...
Cardinality Constraints for Ternary Relations

- Only one single arrow (i.e., cardinality restriction) is allowed for a ternary relation
  - Example: each student can work in at most one project under the supervision of some instructor(s)
Cardinality Constraints for Ternary Relations

- Multiple single arrows (i.e., cardinality restrictions) would lead to different possible interpretations
  - Each student works on at most one project under at most one instructor
  - For each project a student works on, there is at most one instructor
  - For each instructor supervising a student, there is at most one project
- Hence, we do not allow for them
Specialization

- A concept very common in (object oriented) programming
  - Entity sets are sub-/super sets of others
  - They inherit all the attributes from their super sets
- Overlapping
  - A person can be both an employee and a student
- Disjoint
  - An employee can be either an instructor or a secretary
Partial vs. Total Specialization

• Partial specialization
  – An employee may be an instructor or a secretary, or an employee not further specified
  – the default case

• Total specialization
  – There are no other persons than employees and students (in the DB)
  – Needs to be specified in the diagram
  – Analogy in OOP: abstract classes
A Full Example
Reduction to Relation Schemas

• How to get to from an ER model to a relational database model?
  – Recap: relational database models consists of relations

• We have
  – Entity sets and relationship sets

• Goal
  – Translate entity and relationship sets uniformly to relation schemas

• Mechanism:
  – For each entity set and relationship set there is a unique relation that is assigned the name of the corresponding entity set or relationship set
  – Each relation has a number of columns (generally corresponding to attributes), which have unique names
Representing Entity Sets

- A strong entity set becomes a relation with the same attributes
  
  building(name, address)

- A weak entity set becomes a relation that includes
  
  - the column(s) of the primary key of the identifying strong entity set:
    
    room (name, number, capacity)

- At the same time, name is a foreign key
  
  - which integrity constraints should we use?
Representing Relationship Sets

- Many-to-many relationship sets
  - represented as a relation with attributes for the primary keys of the two participating entity sets
- Example: schema for relationship set \( advisor \)

\[
advisor = (\text{student\_ID}, \text{instructor\_ID})
\]

\[
\begin{array}{c|c|c}
\hline
\text{instructor} & \text{advisor} & \text{student} \\
\hline
\text{ID} & & \text{ID} \\
\text{name} & & \text{name} \\
\text{salary} & & \text{tot\_cred} \\
\hline
\end{array}
\]
Representing Relationship Sets

- Many-to-many relationship sets
  - additional attributes of the relationship set become attributes of the representing relation

- Example: schema for relationship set *advisor*

  \[
  \text{advisor} = (\text{student\_ID, instructor\_ID, date})
  \]
Representing Relationship Sets

- Special case for one-to-many relationship sets
  - The primary key of the “many” side can become a foreign key attribute on the “one” side
    \[\text{student} = (\text{ID}, \text{name}, \text{tot\_cred}, \text{instructor\_ID})\]
- In case of partial participation, this may cause null values
Representing Relationship Sets

- Special case for one-to-one relationship sets
  - The primary key on one side can be included on the other side
    
    \[
    \text{student} = (ID, \text{name}, \text{tot_cred}, \text{instructor_ID}) \quad \text{or} \quad \text{instructor} = (ID, \text{name}, \text{salary}, \text{student_id})
    \]

- In case of partial participation, this may cause \text{null} values
Representing Attributes

• Composite attributes are flattened out by creating a separate attribute for each component attribute.

• Add prefix of super attribute in case ambiguous names occur.
  - e.g., street_number, phone_number

• Ignoring multivalued attributes, extended instructor schema is:

  instructor(ID, 
  first_name, middle_initial, last_name, 
  street_number, street_name, 
  apt_number, city, state, zip_code, 
  date_of_birth)
Representing Multi-valued Attributes

- A multivalued attribute $M$ of an entity $E$ is represented by a separate schema $EM$
- Schema $EM$ has attributes corresponding to the primary key of $E$ and an attribute corresponding to multivalued attribute $M$
  - Example: Multivalued attribute $phone\_number$ of instructor is represented by a schema:
    $\text{inst\_phone} = (ID, phone\_number)$
- Each value of the multivalued attribute maps to a separate tuple of the relation on schema $EM$
  - Example: an instructor entity with primary key 22222 and phone numbers 456-7890 and 123-4567 maps to two tuples:
    $(22222, 456-7890)$ and $(22222, 123-4567)$
Representing Derived Attributes

- Technically, we can create a view

```sql
create view instructor_age as
select ID, NOW() - date_of_birth as age
from instructor
```
Representing Higher Arity Relations

- Higher arity relationship sets are represented just like binary ones
  - i.e., as one relation with the primary keys of the related entity sets
  - `proj_guide(instructor_ID, student_ID, project_ID)`
Representing Specialization

• Method 1
  – All three relations become relations
    • primary key is shared
  – Shared attributes are only represented in the higher level entity
    person(ID, name, street, city)
    employee(ID, salary)
    student(ID, tot_credits)

• Drawback:
  – Accessing person information for employees and students requires access to two relations
Representing Specialization

• Method 2
  – All three relations become relations
    • primary key is shared
  – Shared attributes are only represented in each entity
    person(ID, name, street, city)
    employee(ID, name, street, city, salary)
    student(ID, name, street, city, tot_credits)
  – Super relation can be omitted for total specialization

• Drawback:
  – Redundant storage for overlapping specialization
    • i.e., for persons that are both employees and students
Design Decisions in ER Modeling

• Entity sets vs. attributes

  - Entity set
    - Allows for additional information
    - Supports multi-valued attributes
      • in that case, the attribute would end as a relation in the DB anyways

![Diagram showing entity sets and attributes]
Entity Sets vs. Relationship Sets

- Students register for course sections
  - This could be a simple relationship set as well
- Entity set can store additional information, e.g.
  - Date of registration

```
section_reg

registration

student_reg

section

- sec_id
- semester
- year

student

- ID
- name
- tot_cred
```
 Placement of Attributes for 1:1 Relationships

• The primary key on one side can be included on the other side

\[
\text{student} = (\text{ID}, \text{name}, \text{tot_cred}, \text{instructor_ID}) \text{ or } \\
\text{instructor} = (\text{ID}, \text{name}, \text{salary}, \text{student_id})
\]

• Which one?
Binary vs. Non-Binary Relationships

- Sometimes, non-binary relationships can be replaced by binary ones

![Diagram showing person, name, birthday, has_parents, mother, father, and child relationships.](image-url)
Binary vs. Non-Binary Relationships

- Sometimes, non-binary relationships can be replaced by binary ones
  - This is usually the preferred solution
Binary vs. Non-Binary Relationships

- Sometimes, non-binary relationships can be replaced by binary ones
  - but sometimes, they are n-ary by nature
Binary vs. Non-Binary Relationships

- Sometimes, non-binary relationships can be replaced by binary ones
  - but sometimes, they are n-ary by nature
Binary vs. Non-Binary Relationships

• Sometimes, non-binary relationships can be replaced by binary ones
  – but sometimes, they are n-ary by nature

• General decomposition schema:

```
project
ID
name

proj_guide

instructor
ID
name
salary

instr_guide

guide

stud_guide

student
ID
name
tot_credits
```
ER Design Decisions (Summary)

- The use of an attribute or entity set to represent an object
- Whether a real-world concept is best expressed by an entity set or a relationship set
- The use of a ternary relationship versus a number of binary relationships
- The use of a strong or weak entity set
- The use of specialization/generalization – contributes to modularity in the design
Summary of ER Notation

- **E**: entity set
- **R**: relationship set
- **identifying relationship set for weak entity set**
- **total participation of entity set in relationship**

**Attributes**:
- Simple (A1)
- Composite (A2)
- Multivalued (A3)
- Derived (A4)

**Primary Key**
- A1

**Discriminating Attribute**
- A1

Date: 3/9/22
Author: Heiko Paulheim
Summary of ER Notation (ctd.)

- **many-to-many relationship**
- **one-to-one relationship**
- **cardinality limits**
- **ISA: generalization or specialization**
- **disjoint generalization**

Role indicator and role-name are also shown.
Alternative ER Notations

diagram showing:
- entity set E with
  - simple attribute A1,
  - composite attribute A2,
  - multivalued attribute A3,
  - derived attribute A4,
  - and primary key A1

- weak entity set
- generalization
- ISA
- total generalization
Alternative ER Notations (ctd.)

- **many-to-many relationship**
  - E1 * R * E2
  - E1 \(\rightarrow\) R \(\rightarrow\) E2

- **one-to-one relationship**
  - E1 1 R 1 E2
  - E1 \(\rightarrow\) R \(\rightarrow\) E2

- **many-to-one relationship**
  - E1 * R 1 E2
  - E1 \(\rightarrow\) R \(\rightarrow\) E2

- **participation in R: total (E1) and partial (E2)**
  - E1 \(\leftarrow\) R \(\rightarrow\) E2
  - E1 \(\rightarrow\) R \(\rightarrow\) E2
Alternative Modeling Paradigms: UML

• Unified Modeling Language
  – often used in software design
  – similar scope: objects and their relations
  – ISO standard since 2005

• ER models in RDBMS
  – Direct translation to SQL

• UML models in software engineering
  – Direct translation to source code
Alternative Modeling Paradigms: UML

**ER Diagram Notation**

- **E**
  - **A1**: entity with attributes (simple, composite, multivalued, derived)
  - **M10**

- **E1** role1 **R** role2 **E2**: binary relationship

- **E1** role1 **R** role2 **E2**: relationship attributes

- **E1** 0..* **R** 0..1 **E2**: cardinality constraints

**Equivalent in UML**

- **E**
  - **-A1**
  - **+M10**

- **E1** role1 **R** role2 **E2**: class with simple attributes and methods (attribute prefixes: + = public, - = private, # = protected)

- **E1** role1 **R** role2 **E2**: relationship attributes

- **E1** 0..1 **R** 0..* **E2**: cardinality constraints

3/9/22

Heiko Paulheim
Alternative Modeling Paradigms: UML

http://pld.cs.luc.edu/database/ER.html
Summary

• Designing databases
  – i.e., how to get from your customer’s requirements…
  – …to a set of tables and attributes

• ER Models are an intermediate step
  – Conceptual view on the database
  – Graphical notation
  – Can be used for discussion with customers

• Translation rules for ER to RDBMS
• Design decisions
  – For ER Models (mostly business decisions)
  – For translation to RDBMS (mostly computer science decisions)
Questions?