Semantic Web Technologies
Web Ontology Language (OWL)
Part II

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Previously on “Semantic Web Technologies”

• We have got to know
  – OWL, a more powerful ontology language than RDFS
  – Simple ontologies and some reasoning
  – Sudoku solving

• Today
  – New constructs in OWL2
  – Russell's paradox
  – Reasoning in OWL
  – Complexity of ontologies
  – A peek at rule languages for the Semantic Web
Semantic Web – Architecture

here be dragons...

Semantic Web Technologies (This lecture)

Technical Foundations

Berners-Lee (2009): Semantic Web and Linked Data
OWL2 – New Constructs and More

• Five years after the first OWL standard
• OWL2: 2009
  – Syntactic sugar
  – New language constructs
  – OWL profiles

• We have already encountered some, e.g.,
  – Qualified relations
  – Reflexive, irreflexive, and antisymmetric properties
OWL2: Syntactic Sugar

• Disjoint classes and disjoint unions
  – OWL 1:
    
    :Wine owl:equivalentClass [ 
    a owl:Class ; 


  – OWL 2:
    
    :Wine owl:disjointUnionOf 

  – Also possible:
    
    _:x a owl:AllDisjointClasses ;
OWL2: Syntactic Sugar

• Negative(Object|Data)PropertyAssertion
• Allow negated statements
• e.g.: Paul is not Peter's father

  _x [ a owl:NegativeObjectPropertyAssertion;
        owl:sourceIndividual :Paul ;
        owl:targetIndividual :Peter ;

• If that's syntactic sugar, it must also be possible differently
  – But how?
OWL2: Syntactic Sugar

- Negative(Object|Data)PropertyAssertion
- Replaces less intuitive set constructs
- Paul is not Peter's father

```
Paul a [ owl:complementOf [ 
    a owl:Restriction ;
    owl:onProperty :fatherOf ;
    owl:hasValue :Peter
  ];
].
```
OWL2: Reflexive Class Restrictions

- Using `hasSelf`
- Example: defining the set of all autodidacts:

  `:AutoDidact owl:equivalentClass [ a owl:Restriction ; owl:onProperty :teaches ; owl:hasSelf "true"^^xsd:boolean ] .`
OWL2: Profiles

- Profiles are subsets of OWL2 DL
  - EL, RL und QL
  - Similar to complexity classes
- Different runtime and memory complexity
- Depending on requirements
OWL2 Profile

• OWL2 EL (Expressive Language)
  – Fast reasoning on many standard ontologies
  – Restrictions, e.g.:
    • someValuesFrom, but not allValuesFrom
    • No inverse and symmetric properties
    • No unionOf and complementOf

• OWL2 QL (Query Language)
  – Fast query answering on relational databases
  – Restrictions, e.g.:
    • No unionOf, allValuesFrom, hasSelf, …
    • No cardinalities and functional properties
OWL2 Profile

• OWL2 RL (Rule Language)
  – Subset similar to rule languages such as datalog
    • subClassOf is translated to a rule (Person ← Student)
  – Restrictions, e.g.:
    • Only qualified restrictions with 0 or 1
    • Some restrictions for head and body

• The following holds for all three profiles:
  – Reasoning can be implemented in polynomial time for each of the three
  – Reasoning on the union of two profiles only possible in exponential time
OWL2 Example: Russell's Paradox

- A classic paradox by Bertrand Russell, 1918

- In a city, there is exactly one barber who shaves everybody who does not shave themselves.

Who shaves the barber?
OWL2 Example: Russell's Paradox

• Class definitions

  :People owl:disjointUnionOf
  (:PeopleWhoShaveThemselves
   :PeopleWhoDoNotShaveThemselves) .

• Relation definitions:

  :shavedBy rdfs:domain :People .
  :shavedBy rdfs:range :People .
  :shaves owl:inverseOf :shavedBy .

• Every person is shaved by exactly one person:

  :People rdfs:subClassOf [
    a owl:Restriction ;
    owl:onProperty :shavedBy ;
    owl:cardinality "1"^^xsd:integer ] .
OWL2 Example: Russell's Paradox

• Then, we define the barber:

```owl
:Barbers rdfs:subClassOf :People ;
owl:equivalentClass [
    rdf:type owl:Class ;
    owl:oneOf ( :theBarber )
] .
```
OWL2 Example: Russell's Paradox

- Definition of people shaving themselves:

```owl
:PeopleWhoShaveThemselves owl:equivalentClass [
    rdf:type owl:Class ;
    owl:intersectionOf
    ( :People
        [
            a owl:Restriction ;
            owl:onProperty :shavedBy ;
            owl:hasSelf "true"^^xsd:boolean
        ]
    )
] .
```
OWL2 Example: Russell's Paradox

• Definition of people who do not shave themselves:

```owl
:PeopleWhoDoNotShaveThemselves owl:equivalentClass [ a owl:Class ;
owl:intersectionOf ( :People
[ a owl:Restriction
owl:onProperty :shavedBy ;
owl:allValuesFrom :Barbers
]
)
] .
```
OWL2 Example: Russell's Paradox

Help for inconsistent ontologies

Your ontology is inconsistent which means that the OWL reasoner will no longer be able to provide any useful information about the ontology.

You have several options at this point:

- Click the Explain button to try the Protege explanation facility.
- If you think you know what the problem is, click Cancel to fix the ontology yourself.
- Some reasoners come with command line tools that will provide complete explanations for inconsistent ontologies.
OWL2 Example: Russell's Paradox

Inconsistent ontology explanation

1) PersonsWhoDoNotShaveThemselves(?x) -> shaves(the-barber, ?x)
   In 1 other justifications
2) PersonsWhoDoNotShaveThemselves DisjointWith PersonsWhoShaveThemselves
   In ALL other justifications
3) Barber SubClassOf Person
   In ALL other justifications
4) shaves(?x, ?x) -> PersonsWhoShaveThemselves(?x)
   In ALL other justifications
5) shaves(the-barber, ?x) -> PersonsWhoDoNotShaveThemselves(?x)
   In 1 other justifications
6) PersonsWhoShaveThemselves(?x) -> shaves(?x, ?x)
   In ALL other justifications
7) Person EquivalentTo PersonsWhoDoNotShaveThemselves or PersonsWhoShaveThemselves
   In ALL other justifications
8) the-barber Type Barber
   In ALL other justifications

OK
Reasoning in OWL DL

- We have seen reasoning for RDFS
  - Forward chaining algorithm
  - Derive axioms from other axioms

- Reasoning for OWL DL is more difficult
  - Forward chaining may have scalability issues
  - Conjunction (e.g., unionOf) is not supported by forward chaining
  - Different approach: Tableau Reasoning
  - Underlying idea: find contradictions in ontology
    - i.e., both a statement and its opposite can be derived from the ontology
Typical Reasoning Tasks

• What do we want to know from a reasoner?
  – Subclass relations
    • e.g., Are all birds flying animals?
  – Equivalent classes
    • e.g., Are all birds flying animals and vice versa?
  – Disjoint classes
    • e.g., Are there animals that are mammals and birds at the same time?
  – Class consistency
    • e.g., Can there be mammals that lay eggs?
  – Class instantiation
    • e.g., Is Flipper a dolphin?
  – Class enumeration
    • e.g., List all dolphins
Example: A Simple Contradiction

- Given:

:Human a owl:Class .
:Animal a owl:Class .

:Jimmy a :Animal .
:Jimmy a :Human .
Example: A Simple Contradiction

• We can derive:
  – :Human \cap :Animal = \emptyset
    owl:Nothing owl:intersectionOf (:Human :Animal) .
  – :Jimmy \in (:Human \cap :Animal)
    :Jimmy a [ a owl:Class; owl:intersectionOf (:Human :Animal)] .

• i.e.:
  – :Jimmy \in \emptyset
    :Jimmy a owl:Nothing .
  – That means: the instance must not exist
  – but it does
Reasoning Tasks Revisited

• Subclass Relations
  
  Student $\subseteq$ Person $\leftrightarrow$ „Every student is a person“

• Proof method: Reductio ad absurdum
  
  – "Invent" an instance i
  
  – Define Student(i) and $\neg$Person(i)
  
  – Check for contradictions
    
    • If there is one: Student $\subseteq$ Person has to hold
    
    • If there is none: Student $\subseteq$ Person cannot be derived
      
      – Note: it may still hold!
Example: Subclass Relations

• Ontology:
  :Student owl:subClassOf :UniversityMember .
  :UniversityMember owl:subClassOf :Person .

• Invented instance:
  :i a :Student .

• We have
  :i a :Student .
  :Student owl:subClassOf :UniversityMember .

Thus
  :i a :UniversityMember .

• And from
  :UniversityMember owl:subClassOf :Person .

• We further derive that
  :i a Person .
Example: Subclass Relations

• Now, we have

```prolog
:i a :Person .
```

i.e.,

```prolog
:i a [ owl:intersectionOf (:Person [ owl:complementOf :Person ])) .
```

• from which we derive

```prolog
:i a owl:Nothing .
```
Reasoning Tasks Revisited

• Class equivalence
  – Person ≡ Human

• Split into
  – Person ⊆ Human and
  – Human ⊆ Person

• i.e., show subclass relation twice
  – We have seen that

• Class disjointness
  – Are C and D disjoint?
  – "Invent" an instance i
  – Define C(i) and D(i)
    • We have done set (the Jimmy example)
Class Consistency

• Can a class have instances?
  – e.g., married bachelors
    :
    Bachelor owl:subClassOf :Man .
    Bachelor owl:subClassOf
    [ a owl:Restriction;
      owl:onProperty :marriedTo;
      owl:cardinality 0 ] .
    MarriedPerson owl:subClassOf [ a owl:Restriction;
      owl:onProperty :marriedTo;
      owl:cardinality 1 ] .

    MarriedBachelor owl:intersectionOf (:Bachelor :MarriedPerson) .

• Now: invent an instance of the class
  – And check for contradictions
Reasoning Tasks Revisited

• Class Instantiation
  – Is Flipper a dolphin?

• Check:
  – define \neg\text{Dolphin}(\text{Flipper})
  – Check for contradiction

• Class enumeration
  – Repeat class instantiation for all known instances
Typical Reasoning Tasks Revisited

- What do we want to know from a reasoner?
  - Subclass relations
    - e.g., Are all birds flying animals?
  - Equivalent classes
    - e.g., Are all birds flying animals *and vice versa*?
  - Disjoint classes
    - e.g., Are there animals that are mammals and birds at the same time?
  - Class consistency
    - e.g., Can there be mammals that lay eggs?
  - Class instantiation
    - e.g., Is Flipper a dolphin?
  - Class enumeration
    - e.g., List all dolphins
Typical Reasoning Tasks Revisited

• We have seen
  – All reasoning tasks can be reduced to the same basic tasks
  – i.e., consistency checking

• This means: for building a reasoner that can solve those tasks,
  – We only need a reasoner capable of consistency checking
Ontologies in Description Logics Notation

- Classes and Instances
  - $C(x) \leftrightarrow x \text{ a } C$
  - $R(x,y) \leftrightarrow x \mathrel{R} y$
  - $C \sqsubseteq D \leftrightarrow C \text{ rdfs:subClassOf } D$
  - $C \equiv D \leftrightarrow C \text{ owl:equivalentClass } D$
  - $C \sqsubseteq \neg D \leftrightarrow C \text{ owl:disjointWith } D$
  - $C \equiv \neg D \leftrightarrow C \text{ owl:complementOf } D$
  - $C \equiv D \cap E \leftrightarrow C \text{ owl:intersectionOf } (D \ E)$
  - $C \equiv D \cup E \leftrightarrow C \text{ owl:unionOf } (D \ E)$
  - $T \leftrightarrow \text{owl:Thing}$
  - $\bot \leftrightarrow \text{owl:Nothing}$
Ontologies in Description Logics Notation

- **Domains, ranges, and restrictions**
  
  - \( \exists R \cdot T \sqsubseteq C \leftrightarrow R \ rdfs:domain \ C \).
  
  - \( \forall R \cdot C \leftrightarrow R \ rdfs:range \ C \).
  
  - \( C \sqsubseteq \forall R \cdot D \leftrightarrow C \ owl:subClassOf \)
    \[
    [ \ a \ owl:Restriction;
      \ owl:onProperty \ R;
      \ owl:allValuesFrom \ D \ ] .
    \]
  
  - \( C \sqsubseteq \exists R \cdot D \leftrightarrow C \ owl:subClassOf \)
    \[
    [ \ a \ owl:Restriction;
      \ owl:onProperty \ R;
      \ owl:someValuesFrom \ D \ ] .
    \]
  
  - \( C \sqsubseteq \geq n R \leftrightarrow C \ owl:subClassOf \)
    \[
    [ \ a \ owl:Restriction;
      \ owl:onProperty \ R;
      \ owl:minCardinality \ n \ ] .
    \]
Global Statements in Description Logic

- So far, we have seen mostly statements about single classes
  - e.g., $C \sqsubseteq D$

- In Description Logics, we can also make global statements
  - e.g., $D \sqcup E$
  - This means: every single instance is a member of $D$ or $E$ (or both)

- Those global statements are heavily used in the reasoning process
Negation Normal Form (NNF)

- Transforming ontologies to Negation Normal Form:
  - $\exists$ und $\equiv$ are not used
  - Negation only for atomic classes and axioms

- A simplified notation of ontologies
- Used by tableau reasoners
Negation Normal Form (NNF)

• Eliminating $\equiv$:
  • Replace $C \equiv D$ by $\neg C \lor D$
  • Note: this is a shorthand notation for $\forall x: \neg C(x) \lor D(x)$

• Why does this hold?
  • $C \equiv D$ is equivalent to $C(x) \rightarrow D(x)$

<table>
<thead>
<tr>
<th>C(x)</th>
<th>D(x)</th>
<th>C(x) → D(x)</th>
<th>$\neg C(x) \lor D(x)$</th>
</tr>
</thead>
<tbody>
<tr>
<td>true</td>
<td>true</td>
<td>true</td>
<td>true</td>
</tr>
<tr>
<td>true</td>
<td>false</td>
<td>false</td>
<td>false</td>
</tr>
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</tr>
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<td>false</td>
<td>false</td>
<td>true</td>
<td>true</td>
</tr>
</tbody>
</table>
Negation Normal Form (NNF)

• Eliminating $\equiv$:
  • Replace $C \equiv D$ by $C \sqsubseteq D$ and $D \sqsubseteq C$
  • Proceed as before

• i.e.: $C \equiv D$ becomes

  $C \sqsubseteq D$
  $D \sqsubseteq C$

  – and thus

  $\neg C \sqcup D$
  $\neg D \sqcup C$
Negation Normal Form (NNF)

- Further transformation rules
  - $\text{NNF}(C) = C$ (for atomic $C$)
  - $\text{NNF}(\neg C) = \neg C$ (for atomic $C$)
  - $\text{NNF}(\neg \neg C) = C$
  - $\text{NNF}(C \lor D) = \text{NNF}(C) \lor \text{NNF}(D)$
  - $\text{NNF}(C \land D) = \text{NNF}(C) \land \text{NNF}(D)$
  - $\text{NNF}(\neg (C \land D)) = \text{NNF}(\neg C) \lor \text{NNF}(\neg D)$
  - $\text{NNF}(\neg (C \lor D)) = \text{NNF}(\neg C) \land \text{NNF}(\neg D)$
  - $\text{NNF}(\forall R.C) = \forall R.\text{NNF}(C)$
  - $\text{NNF}(\exists R.C) = \exists R.\text{NNF}(C)$
  - $\text{NNF}(\neg \forall R.C) = \exists R.\text{NNF}(\neg C)$
  - $\text{NNF}(\neg \exists R.C) = \forall R.\text{NNF}(\neg C)$
The Basic Tableau Algorithm

• Tableau: Collection of derived axioms
  – Is subsequently extended
  – As for forward chaining

• In case of conjunction
  – Split the tableau

<table>
<thead>
<tr>
<th>C(a)</th>
<th>D(a) ∪ E(a)</th>
<th>C(a), D(a)</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td>C(a), E(a)</td>
</tr>
</tbody>
</table>
When is an Ontology Free of Contradictions?

• Tableau is continuously extended and split
• Free of contradictions if...
  – No further axioms can be created
  – At least one partial tableau is free of contradictions
  – A partial tableau has a contradiction if it contains both an axiom and its negation
    • e.g., Person(Peter) und ¬Person(Peter)
    • The partial tableau is then called closed
The Basic Tableau Algorithm

- Given: an ontology O in NNF

  While not all partial tableaus are closed
    * Choose a non-closed partial tableau T and an $A \in O \cup T$
      If A is not contained in T
        If A is an atomic statement
          add A to T
          back to *
        If A is a non-atomic statement
          Choose an individual $i \in O \cup T$
          Add A(i) to T
          back to *
      else
        Extend the tableau with consequences from A
        back to *
The Basic Tableau Algorithm

- Extending a tableau with consequences

<table>
<thead>
<tr>
<th>Nr</th>
<th>Axiom</th>
<th>Action</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>C(a)</td>
<td>Add C(a)</td>
</tr>
<tr>
<td>2</td>
<td>R(a,b)</td>
<td>Add R(a,b)</td>
</tr>
<tr>
<td>3</td>
<td>C</td>
<td>Choose an individual a, add C(a)</td>
</tr>
<tr>
<td>4</td>
<td>(C \sqcap D)(a)</td>
<td>Add C(a) and D(a)</td>
</tr>
<tr>
<td>5</td>
<td>(C \sqcup D)(a)</td>
<td>Split tableau into T1 and T2. Add C(a) to T1, D(a) to T2</td>
</tr>
<tr>
<td>6</td>
<td>(\exists R.C)(a)</td>
<td>Add R(a,b) and C(b) for a new Individual b</td>
</tr>
<tr>
<td>7</td>
<td>(\forall R.C)(a)</td>
<td>For all b with R(a,b) \in T: add C(b)</td>
</tr>
</tbody>
</table>
A Simple Example

• Given the following ontology:

  :Seth a :Human.
  :Seth a :Insect.

• Is this knowledge base consistent?
A Simple Example

• Given the following ontology:
  :Seth a :Human .
  :Seth a :Insect .

  – The same ontology in DL-NNF:
    ¬Animal ⊔ ¬Human
    Animal ⊔ (¬Mammal ⊓ ¬Bird ⊓ ¬Fish ⊓ ¬Insect ⊓ ¬Reptile)
    ¬Animal ⊔ (Mammal ⊔ Bird ⊔ Fish ⊔ Insect ⊔ Reptile)
    Human(Seth)
    Insect(Seth)

• Let's try how reasoning works now!
A Simple Example

Human(Seth), Insect(Seth)

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<tbody>
<tr>
<td>1</td>
<td>C(a)</td>
<td>Add C(a)</td>
</tr>
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</table>
A Simple Example

Human(Seth), Insect(Seth),
(¬Animal ⊔ ¬Human)(Seth)

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</tr>
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<tbody>
<tr>
<td>3</td>
<td>C</td>
<td>Choose an individual a, add C(a)</td>
</tr>
</tbody>
</table>
A Simple Example

Human(Seth), Insect(Seth),
¬Animal(Seth)

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<tr>
<td>5</td>
<td>(C [\lor] D)(a)</td>
<td>Split the tableau into T1 and T2. Add C(a) to T1, D(a) to T2</td>
</tr>
</tbody>
</table>
**A Simple Example**

\[
\text{Human(Seth), Insect(Seth),} \\
\neg \text{Animal(Seth)} \\
\text{Animal} \cup (\neg \text{Mammal} \land \neg \text{Bird} \land \neg \text{Fish} \land \neg \text{Insect})(\text{Seth})
\]

\[
\text{Human(Seth), Insect(Seth),} \\
\neg \text{Human(Seth)}
\]

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A Simple Example

Human(Seth), Insect(Seth),
¬Animal(Seth)
Animal(Seth)

Human(Seth), Insect(Seth),
¬Animal(Seth)
(¬Mammal □ ¬Bird □ ¬Fish □ ¬Insect □ ¬Reptile)(Seth)

Human(Seth), Insect(Seth),
¬Human(Seth)

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### A Simple Example

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<td>4</td>
<td>((C \sqcap D)(a))</td>
<td>Add (C(a)) and (D(a))</td>
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Another Example

- Again, a simple ontology:

```prolog
:Woman rdfs:subClassOf :Person .
:Man rdfs:subClassOf :Person .
:hasChild rdfs:domain :Person .
:hasChild rdfs:range :Person .
:Peter :hasChild :Julia .
:Julia a :Woman .
:Peter a :Man .
```
Another Example

• in DL NNF:

\neg \text{Man} \sqsubseteq \text{Person}
\neg \text{Woman} \sqsubseteq \text{Person}
\exists \text{hasChild} \sqsubseteq \text{Person}
\forall \text{hasChild} \cdot \text{Person}
\text{hasChild(Peter,Julia)}
\text{Woman(Julia)}
\text{Man(Peter)}
Another Example

hasChild(Peter, Julia)

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<td>2</td>
<td>R(a,b)</td>
<td>Add R(a,b)</td>
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Another Example

hasChild(Peter, Julia), Woman(Julia)

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<tbody>
<tr>
<td>1</td>
<td>C(a)</td>
<td>Add C(a)</td>
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</table>
Another Example

\[
\text{hasChild}(\text{Peter}, \text{Julia}), \text{Woman}(\text{Julia}), \\
(\neg \exists \text{hasChild}. T \sqsubseteq \text{Person})(\text{Peter})
\]

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hasChild(Peter, Julia), Woman(Julia),
(¬∃hasChild.T ⊔ Person)(Peter),
¬∃hasChild.T(Peter)

hasChild(Peter, Julia), Woman(Julia),
(¬∃hasChild.T)(Peter),
Person(Peter)
Another Example

\[
\text{hasChild}(\text{Peter}, \text{Julia}), \text{Woman(Julia)}, \\
(\neg \exists \text{hasChild}.T \sqcup \text{Person})(\text{Peter}), \\
\neg \exists \text{hasChild}.T(\text{Peter})
\]

\[
\text{hasChild}(\text{Peter}, \text{Julia}), \text{Woman(Julia)}, \\
(\neg \exists \text{hasChild}.T)(\text{Peter}), \\
\text{Person}(\text{Peter}), \\
\neg \text{hasChild}(\text{Peter}, \text{b0}), \text{T(b0)}
\]

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<td>6</td>
<td>$(\exists R.C)(a)$</td>
<td>Add $R(a,b)$ und $C(b)$ for a <em>new</em> Individual $b$</td>
</tr>
</tbody>
</table>
Another Example

\[
\begin{align*}
\text{hasChild(Peter, Julia), Woman(Julia),} \\
(\neg \exists \text{hasChild.T } \sqcap \text{Person})(Peter), \\
\neg \exists \text{hasChild.T}(Peter)
\end{align*}
\]

\[
\begin{align*}
\text{hasChild(Peter, Julia), Woman(Julia),} \\
(\neg \exists \text{hasChild.T})(Peter), \\
\text{Person(Peter),} \\
\neg \text{hasChild(Peter,b0),T(b0),} \\
\neg \text{hasChild(Peter,b1),T(b1),} \\
\text{...}
\end{align*}
\]

<table>
<thead>
<tr>
<th>Nr</th>
<th>Axiom</th>
<th>Action</th>
</tr>
</thead>
<tbody>
<tr>
<td>6</td>
<td>(\exists R.C)(a)</td>
<td>Add R(a,b) und C(b) for a new Individual b</td>
</tr>
</tbody>
</table>
Introducing Rule Blocking

• Observation
  – The tableau algorithm does not necessarily terminate
  – We can add arbitrarily many new axioms

<table>
<thead>
<tr>
<th>Nr</th>
<th>Axiom</th>
<th>Action</th>
</tr>
</thead>
<tbody>
<tr>
<td>6</td>
<td>(∃R.C)(a)</td>
<td>Add R(a,b) und C(b) for a new Individual b</td>
</tr>
</tbody>
</table>

• Idea: avoid rule 6 if no new information is created
  – i.e., if we already created one instance \( b_a \) for instance \( a \),
    then block using rule 6 for \( a \).
Tableau Algorithm with Rule Blocking

- **Given:** an ontology \( O \) in NNF

  While not all partial tableaus are closed and further axioms can be created

  * Choose a non-closed partial tableau \( T \) and a non-blocked \( A \in O \cup T \)
    
    If \( A \) is not contained in \( T \)
    
    If \( A \) is an atomic statement
      
      add \( A \) to \( T \)
      
      back to *
    
    If \( A \) is a non-atomic statement
      
      Choose an individual \( i \in O \cup T \)
      
      Add \( A(i) \) to \( T \)
      
      back to *
    
    else
    
    Extend the tableau with consequences from \( A \)
    
    If rule 6 was used, block \( A \) for \( T \)
    
    back to *
Tableau Algorithm: Wrap Up

- An algorithm for description logic based ontologies
  - works for OWL Lite and DL
- We have seen examples for some OWL expressions
  - Other OWL DL expressions can be “translated” to DL as well
  - And they come with their own expansion rules
  - Reasoning may become more difficult
    - e.g., dynamic blocking and unblocking
Optimizing Tableau Reasoners

• Given: an ontology O in NNF

While not all partial tableaus are closed and further axioms can be created

* Choose a non-closed partial tableau T and a non-blocked A ∈ O ∪ T

  If A is not contained in T

    If A is an atomic statement

      add A to T
      back to *

    If A is a non-atomic statement

      Choose an individual i ∈ O ∪ T
      Add A(i) to T
      back to *

  else

    Extend the tableau with consequences from A

  If rule 6 was used, block A for T

  back to *
OWL Lite vs DL Revisited

• Recap: OWL Lite has some restrictions
  – Those are meant to allow for faster reasoning

• Restrictions only with cardinalities 0 and 1
  – Higher cardinalities make blocking more complex

• unionOf, disjointWith, complementOf, closed classes, ...
  – they all introduce more disjunctions
  – i.e., more splitting operations
Complexity of Ontologies

• Reasoning is usually expensive

• Reasoning performance depends on ontology complexity
  – Rule of thumb: the more complexity, the more costly

• Most useful ontologies are in OWL DL
  – But there are differences
  – In detail: complexity classes
Simple Ontologies: ALC

• ALC: Attribute Language with Complement

• Allowed:
  – subClassOf, equivalentClass
  – unionOf, complementOf, disjointWith
  – Restrictions: allValuesFrom, someValuesFrom
  – domain, range
  – Definition of individuals
 Complexity classes are noted as letter sequences

Using
  - S = ALC plus transitive properties (basis for most ontologies)
  - H = Property hierarchies (subPropertyOf)
  - O = closed classes (oneOf)
  - I = inverse properties (inversePropertyOf)
  - N = numeric restrictions (min/maxCardinality)
  - F = functional properties
  - Q = qualified numerical restrictions (OWL2)
  - (D) = Usage of datatype properties
Some Tableau Reasoners

• Fact
  – University of Manchester, free
  – SHIQ

• Fact++/JFact
  – Extension of Fact, free
  – SHOIQ(and a little D), OWL-DL + OWL2

• Pellet
  – Clark & Parsia, free for academic use
  – SHOIN(D), OWL-DL + OWL2

• RacerPro
  – Racer Systems, commercial
  – SHIQ(D)
Sudoku Revisited

• Recap: we used a closed class
  – Plus some disjointness
• Resulting complexity: SO
• Which reasoners do support that?
  – Fact: SHIQ :-(
  – RacerPro: SHIQ(D) :-(
  – Pellet: SHOIN(D) :-)
  – HermiT: SHOIQ :-)

\[
\begin{array}{ccc|ccc|ccc}
5 & 3 & & 7 & & & & & \\
6 & & 1 & 9 & 5 & & & & \\
 & 9 & 8 & & & 6 & & & \\
8 & & 6 & & 3 & & 1 & & \\
4 & 8 & 3 & & 1 & & & & \\
7 & & 2 & & 6 & & & & \\
 & 6 & & & & 2 & & 8 & \\
 & 4 & 1 & 9 & & 5 & & & \\
 & 8 & & 7 & 9 & & & & \\
\end{array}
\]
Rules: Beyond OWL

here be dragons...

Semantic Web Technologies
(This lecture)

Technical Foundations

Berners-Lee (2009): Semantic Web and Linked Data
Limitations of OWL

• Some things are hard or impossible to express in OWL
• Example:
  – If A is a woman and the child of B
    then A is the daughter of B
Limitations of OWL

- Let's try this in OWL:

```owl
:Woman rdfs:subClassOf :Human .
:childOf a owl:ObjectProperty ;
  rdfs:domain :Human ;
  rdfs:range :Human .
:daughterOf a owl:ObjectProperty ;
  rdfs:subPropertyOf :childOf ;
  rdfs:domain :Woman .
```
Limitations of OWL

• What can a reasoner conclude with this ontology?
• Example:

  :Julia :daughterOf :Peter .
  → :Julia a :Woman .

• What we would like to have instead:

  :Julia :childOf :Peter .
  :Julia a :Woman .
  → :Julia :daughterOf :Peter .
Limitations of OWL

• What we would like to have:
  \[ \text{daughterOf}(X,Y) \leftarrow \text{childOf}(X,Y) \land \text{Woman}(X). \]

• Rules are flexible
• There are rules in the Semantic Web, e.g.
  – Semantic Web Rule Language (SWRL)
  – Rule Interchange Format (RIF)
  – See lecture in two weeks
• Some reasoners do (partly) support rules
Wrap Up

• OWL comes in many flavours
  – OWL Lite, OWL DL, OWL Full
  – Detailed complexity classes of OWL DL
  – Additions and profiles from OWL2
  – However, there are still some things that cannot be expressed...

• Reasoning is typically done using the Tableau algorithm
Questions?